

CS310

Chomsky Normal Form

Section: 2.1

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Pacific University

Quick Review

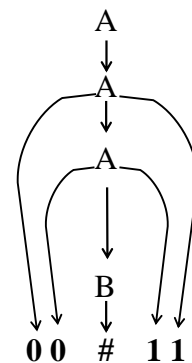
- (CFG) 4-tuple (V, Σ, R, S)
- V finite set of variables
- Σ finite set of terminals
- R set of rules of form:
 - variable \rightarrow (string of variables and terminals)
- $S \in V$, start variable
- $L(G) = \{ w \in \Sigma^* \mid S \rightarrow^* w \}$
- w is in Σ^* and can be derived from S

Example

$A \rightarrow 0A1$

$A \rightarrow B$

$B \rightarrow \#$



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Chomsky Normal Form

•CNF presents a grammar in a standard, simplified form:

$$A \rightarrow BC$$

$$A \rightarrow a$$

1. where A,B,C are variables and B and C are not the start variable AND
2. a is a terminal

NOTE: The rule $S \rightarrow \epsilon$ (S Start Variable) is also allowed so the language can generate the empty string (optional)

Chomsky Normal Form Examples

Which CFG's are in Chomsky Normal Form?

- $S \rightarrow XM$
 $M \rightarrow SY$
 $X \rightarrow x$
 $Y \rightarrow y$
- $S \rightarrow xSy$
 $S \rightarrow xy$

CNF Benefits

- Easier to prove statements about CFG's when in CNF
- Any CFG can be converted to CNF
- Remove productions:
 1. $A \rightarrow \epsilon$ to empty
 2. $A \rightarrow B$ Unit rule
 3. $A \rightarrow s$, s contains a terminal and $|s| > 1$
 4. $A \rightarrow s$, $|s| > 2$, $s \in \{ V U \Sigma \}^*$

Step 1: Removing $A \rightarrow \epsilon$

Example:

$S \rightarrow UAV$

$A \rightarrow \epsilon$

In CNF? If not, how to fix?

Nullable variables

- A variable A is *nullable* if $A \rightarrow^* \epsilon$

So we must:

- Find all nullable variables
- Remove all ϵ transitions
- If $T \rightarrow X_1AX_2$ and A is nullable then add $T \rightarrow X_1X_2$

Step 1 Example

$S \rightarrow TU$

$T \rightarrow AB$

$A \rightarrow aA \mid \epsilon$

$B \rightarrow bB \mid \epsilon$

$U \rightarrow ccA \mid B$

Nullable variables?

Productions removed?

Productions added?

Step 2: Removing $A \rightarrow B$ (Unit Productions)

$A \rightarrow B$ Must add $A \rightarrow s$ to compensate!

$B \rightarrow s$

$S \in \{ V \cup \Sigma \}^*$

• A variable B is **A-derivable** if $A \rightarrow^* B$

• Must:

- Find all A-derivable variables for each A
- Remove all unit transitions
- If $B \rightarrow s$ and B is A-derivable then add $A \rightarrow s$

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Step 2 Example

$S \rightarrow TU \mid T \mid U$ $B \rightarrow bB \mid b$

$T \rightarrow AB \mid A \mid B$ $U \rightarrow ccA \mid B \mid cc$

$A \rightarrow aA \mid a$

S-derivable:

T-derivable:

U-derivable:

NOTE: Not concerned about A, B since not
involved in unit production

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Step 2 Example, cont.

$S \rightarrow TU \mid T \mid U$ $B \rightarrow bB \mid b$
 $T \rightarrow AB \mid A \mid B$ $U \rightarrow ccA \mid B \mid cc$
 $A \rightarrow aA \mid a$

Productions removed:

Productions added:

Step 2 Example, cont.

$S \rightarrow TU \mid T \mid U \mid \varepsilon$ $B \rightarrow bB \mid b$
 $T \rightarrow AB \mid A \mid B$ $U \rightarrow ccA \mid B \mid cc$
 $A \rightarrow aA \mid a$

Productions that still violate CNF?:

Steps 3,4 Example, cont.

Grammar should now have NO unit productions or null productions $A \rightarrow \epsilon$ (except A may be start variable in null production)

Still need to fix productions with terminals on the RHS

Steps 3, 4: Remove $A \rightarrow S_1 a S_2$

$$A \rightarrow S_1 a S_2$$

$a \in \Sigma$, S_1 and S_2 are strings, at least one is not empty

Create

$$X_a \rightarrow a$$

$$A \rightarrow S_1 X_a S_2$$

Then fix up $A \rightarrow S_1 X_a S_2$

–why? what rule is violated?

–how?

Remove $A \rightarrow S_1 X_a S_2$

$A \rightarrow S_1 X_a S_2$

$A \rightarrow$

Note: In general, chain productions:

$A \rightarrow A_1 A_2 \dots A_k$, then decompose as:

$A \rightarrow A_1 W_1, W_1 \rightarrow A_2 W_2 \dots W_{k-1} \rightarrow A_{k-1} A_k$

Step 3,4 Example, cont.

$S \rightarrow TU \mid aA \mid A \mid bB \mid AB \mid ccA \mid cc \mid \varepsilon$

$T \rightarrow AB \mid aA \mid a \mid bB \mid b$

$A \rightarrow aA \mid a$

$B \rightarrow bB \mid b$

$U \rightarrow ccA \mid cc \setminus bB \setminus b$

$A \rightarrow aA \mid a$

Finish!

$S \rightarrow ASA \mid aB$

Put in to CNF

$A \rightarrow B \mid S$

$B \rightarrow b \mid \varepsilon$

Should get a grammar equivalent to:

$S_0 \rightarrow AA_1 \mid X_a B \mid a \mid SA \mid AS$

$S \rightarrow AA_1 \mid X_a B \mid a \mid SA \mid AS$

$A \rightarrow b \mid a A_1 \mid X_a B \mid a \mid SA \mid AS$

$A_1 \rightarrow SA$

$X_a \rightarrow a$

$B \rightarrow b$