CS480
Syntax Directed Translation
Extra Slides

March 13, 2013

Semantic Rules

- Problems?

```
int main()
{
  x = 9;
  return 0;
}
```

```
int main()
{
  int y, x;
  y = x;
  return 0;
}
```

```
int main()
{
  int y, x;
  if ( input() == 1 )
  {
    x = 0;
    y = x;
  }
  y = x;
  return 0;
}
```

Data Flow Analysis p 608

- Separate from compilers
  - other use of a parser

  - think of a parse tree based on statements rather than tokens
  - imagine we replace ***statement*** in our grammar like we did EXPRESSION in TD

  – reaching definitions

  – live variables

  – definition-use chain

  – Equation:
    - out[S] = gen[S] U (in[S] - kill[S])

```
exp -> lvalue assignop exp | deref assignop exp | orterm
deref -> * unexp
orterm -> orterm || andterm | andterm
andterm -> andterm && eqterm | eqterm
eqterm -> eqterm equop relterm | relterm
relterm -> relterm relop term | term
term -> term addop factor | factor
factor -> factor mulop unexp | unexp
unexp -> lvalue autoop | & lvalue | * unexp | negop unexp
        | primary
primary -> ( exp ) | lvalue | constant | func
lvalue -> var | ( lvalue )
var -> id | id [ exp ]
func -> id | id ( list )
list -> exp | list , exp
```

Example

- How does the compiler handle ***overloaded*** functions?
  - why can you only do this in C++/Java not C?
  - what does **`extern "C"`** mean?

```
void foo(x,y)
int x;
int y;
{ }
```

```
void foo(x,y)
int x;
int *y;
{ }
```

Semantic Rules?

I=int, P=ptr, A*X*=array (*X*=I or P)
foo_II foo_IP

```
externaldefs -> int externaldef
              | ε
externaldef  -> id typepart externaldefs
              | * id vardecl ; externaldefs
typepart     -> ( optparamlist ) functionbody
              | vardecl ;
functionbody -> typedecllist functionstmt
typedecllist -> int idorptr optarraydecl moredecls ; typedecllist
              | ε
idorptr      -> id
              | * id
optarraydecl -> [ ]
              | ε
moredecls    -> , idorptr optarraydecl moredecls
              | ε
```