
CS480

Syntax Directed Translation

Extra Slides

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Semantic Rules

- Problems?

```
int main()
{
    x = 9;
    return 0;
}
```

```
int main()
{
    int y, x;
    y = x;
    return 0;
}
```

```
int main()
{
    int y, x;
    if ( input() == 1 )
    {
        x = 0;
        y = x;
    }
    y = x;
    return 0;
}
```

Data Flow Analysis p 608

- Separate from compilers
 - other use of a parser
- think of a parse tree based on statements rather than tokens
- imagine we replace *statement* in our grammar like we did EXPRESSION in TD
- reaching definitions
- live variables
- definition-use chain
- Equation:
 - $out[S] = gen[S] \cup (in[S] - kill[S])$

```

exp -> lvalue assignop exp | * unexp assignop exp | orterm
orterm -> orterm || andterm | andterm
andterm -> andterm && eqterm | eqterm
eqterm -> eqterm equop relterm | relterm
relterm -> relterm relop term | term
term -> term addop factor | factor
factor -> factor mulop unexp | unexp
unexp -> lvalue autoop | & lvalue | * unexp | negop unexp
      | primary
primary -> ( exp ) | lvalue | constant | func
lvalue -> var | ( lvalue )
var -> id | id [ exp ]
func -> id | id ( list )
list -> exp | list , exp

```

Example

- How does the compiler handle *overloaded* functions?
 - why can you only do this in C++/Java not C?
 - what does **extern** "C" mean?

```
void foo(x,y)
int x;
int y;
{ }
```

```
void foo(x,y)
int x;
int *y;
{ }
```

Semantic Rules?

I=int, P=ptr, AX=array (X=I or P)

foo_II foo_IP

```
externaldefs -> int externaldef
              |  $\epsilon$ 
externaldef  -> id typepart externaldefs
              | * id vardecl ; externaldefs
typepart     -> ( optparamlist ) functionbody
              | vardecl ;
functionbody -> typedecllist functionstmt
typedecllist -> int idorptr optarraydecl moredecls ; typedecllist
              |  $\epsilon$ 
idorptr      -> id
              | * id
optarraydecl -> [ ]
              |  $\epsilon$ 
moredecls    -> , idorptr optarraydecl moredecls
              |  $\epsilon$ 
```