
Red-Black Trees

Chapters 13

Binary Search Trees Review

Balanced Trees

- Why do we want to balance trees?
- Red-Black Trees are an example of balanced trees
- Other balanced trees:
 - AVL trees
 - B-trees
 - 2-3 trees

Red-Black Tree

- BST data structure with extra color field for each node, satisfying the red-black properties:
 1. Every node is either red or black.
 2. The root is black.
 3. Every leaf is black.
 4. If a node is red, both children are black.
 5. Every path from node to descendent leaf contain the same number of black nodes.

Example

- Attributes of nodes:
 - key
 - left
 - right
 - p (parent)
 - color
- Note the use of the sentinel T.nil
 - Parent of the root is T.nil
 - All leaves are T.nil – no data in the leaves!

Properties of RB-Trees

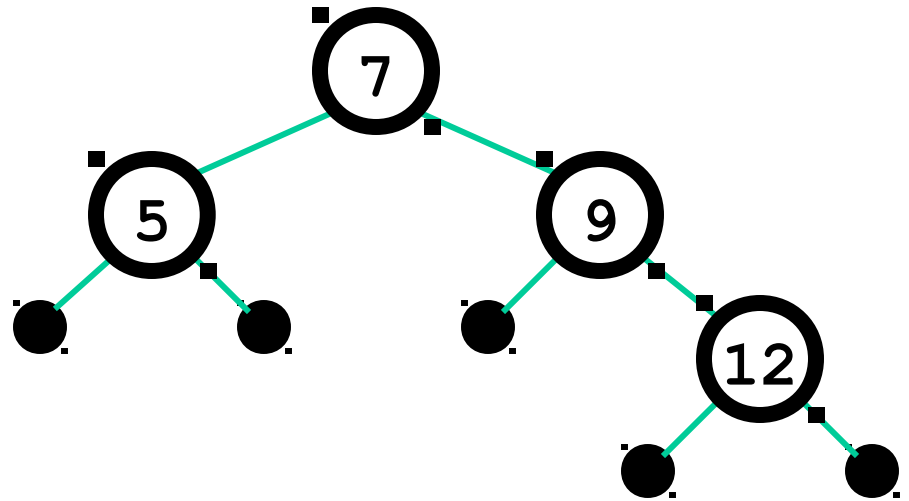
- Black-height of a node:
 - Number of black nodes on any simple path from, but not including, a node x down to a leaf
- A red-black tree with n internal nodes has height at most $2\lg(n+1)$

Rotations

- Why are rotations necessary in red-black trees?
- How are rotations performed?
- What is the running time of rotation?

Example

- Color this tree
- Insert 8
- Insert 11
- Insert 10



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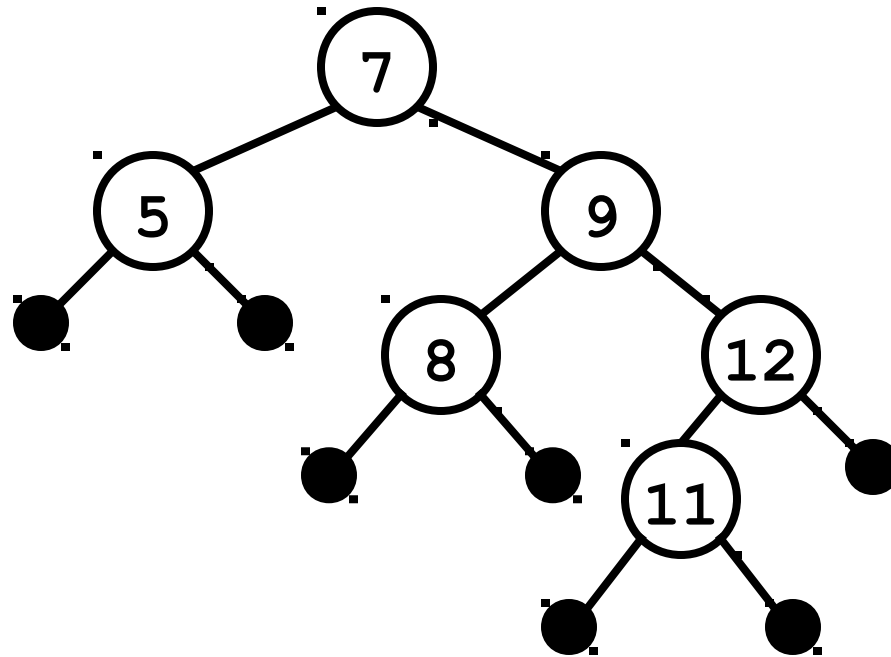
Left-Rotate

LEFT-ROTATE(T, x)

```
y = x.right           // set y
x.right = y.left     // turn y's left subtree into x's right subtree
if y.left ≠ T.nil
    y.left.p = x
y.p = x.p           // link x's parent to y
if x.p == T.nil
    T.root = y
elseif x == x.p.left
    x.p.left = y
else x.p.right = y
y.left = x         // put x on y's left
x.p = y
```

Example

- Rotate left about 9



Inserting into a RB-Tree

- This is regular binary search tree insertion
- Which RB-Tree property could have been violated?

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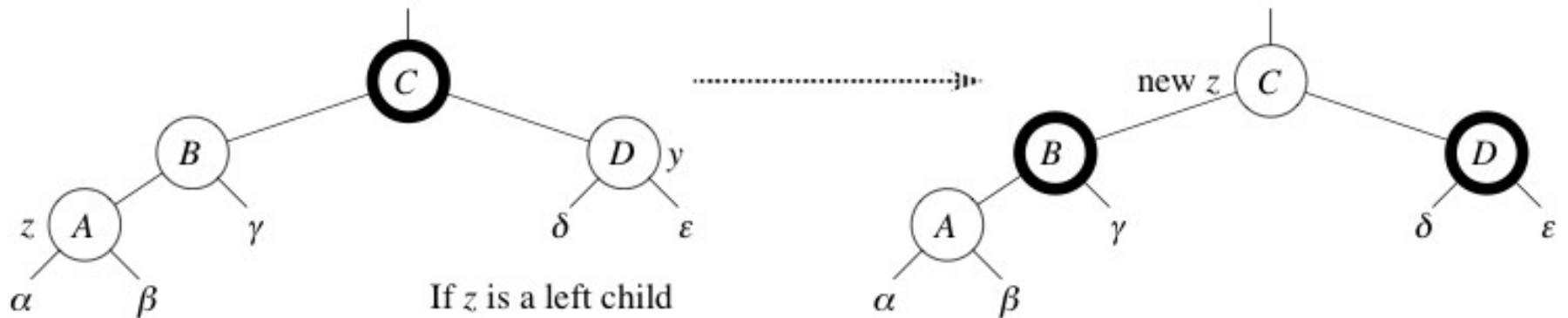
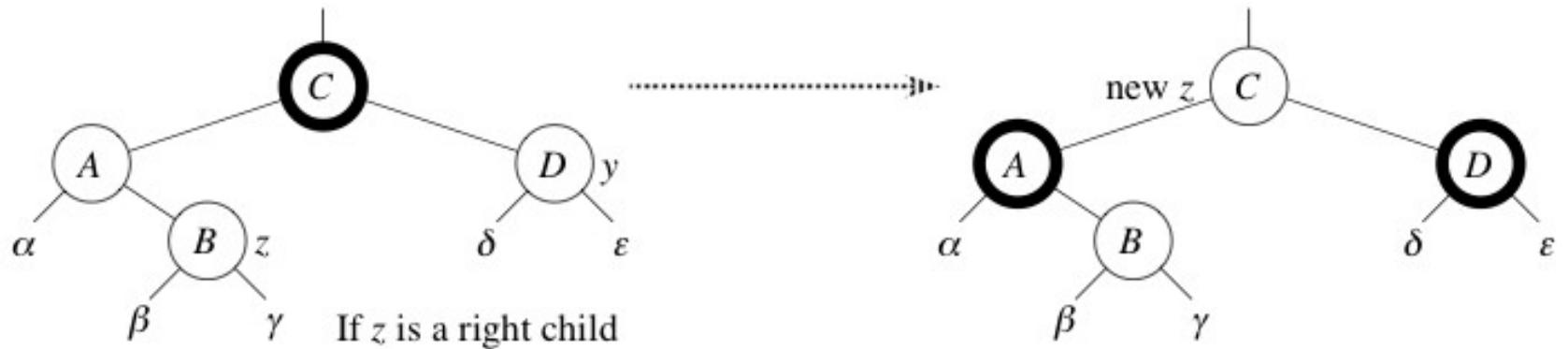
```
RB-INSERT(T, z)
  y = T.nil
  x = T.root
  while x ≠ T.nil
    y = x
    if z.key < x.key
      x = x.left
    else x = x.right
  z.p = y
  if y == T.nil
    T.root = z
  elseif z.key < y.key
    y.left = z
  else y.right = z
  z.left = T.nil
  z.right = T.nil
  z.color = RED
  RB-INSERT-FIXUP(T, z)
```

RB-Insert-Fixup

```
RB-INSERT-FIXUP(T, z)
while z.p.color == RED
    if z.p == z.p.p.left
        y = z.p.p.right
        if y.color == RED
            z.p.color = BLACK // case 1
            y.color = BLACK // case 1
            z.p.p.color = RED // case 1
            z = z.p.p // case 1
        else if z == z.p.right
            z = z.p // case 2
            LEFT-ROTATE(T, z) // case 2
            z.p.color = BLACK // case 3
            z.p.p.color = RED // case 3
            RIGHT-ROTATE(T, z.p.p) // case 3
        else (same as then clause with “right” and “left” exchanged)
    T.root.color = BLACK
```

Cases

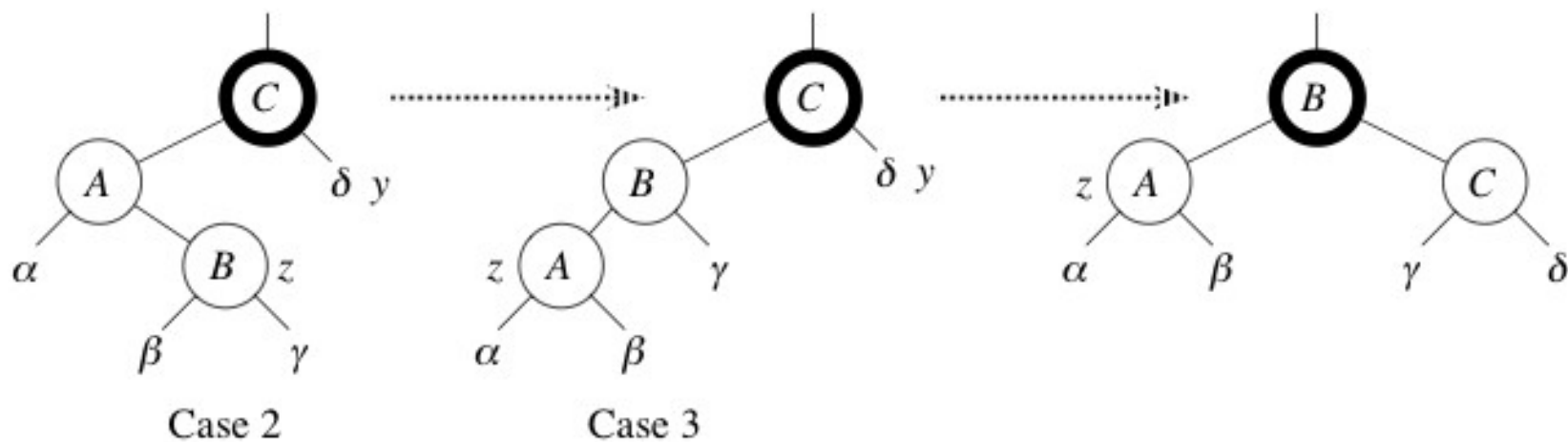
Case 1: y is red



Cases

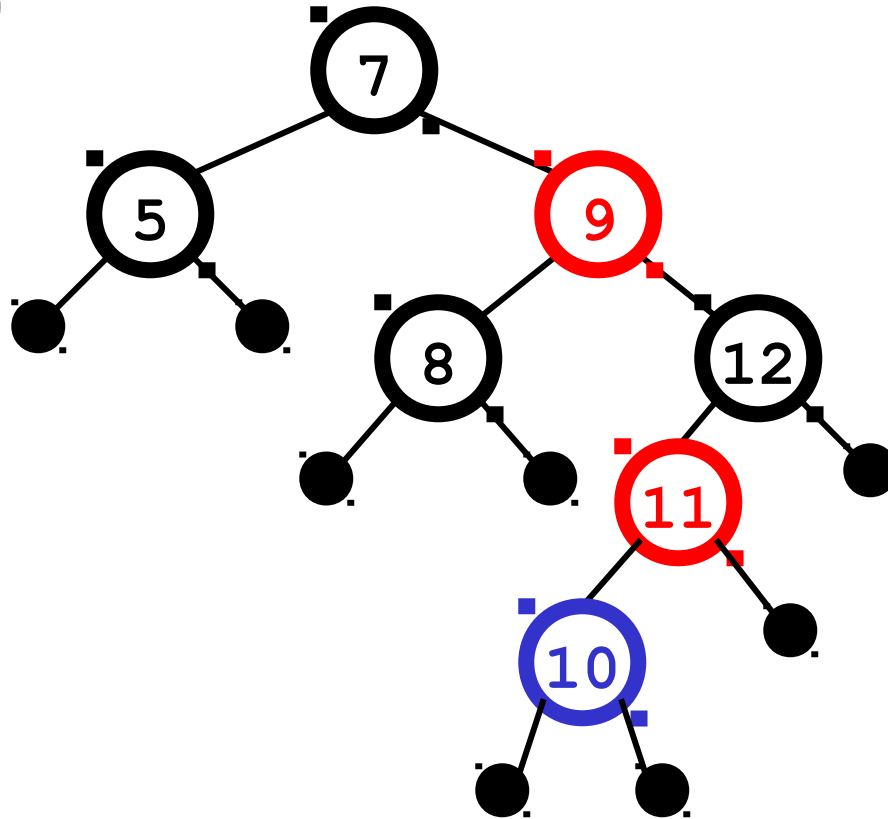
Case 2: y is black, z is a right child

Case 3: y is black, z is a left child



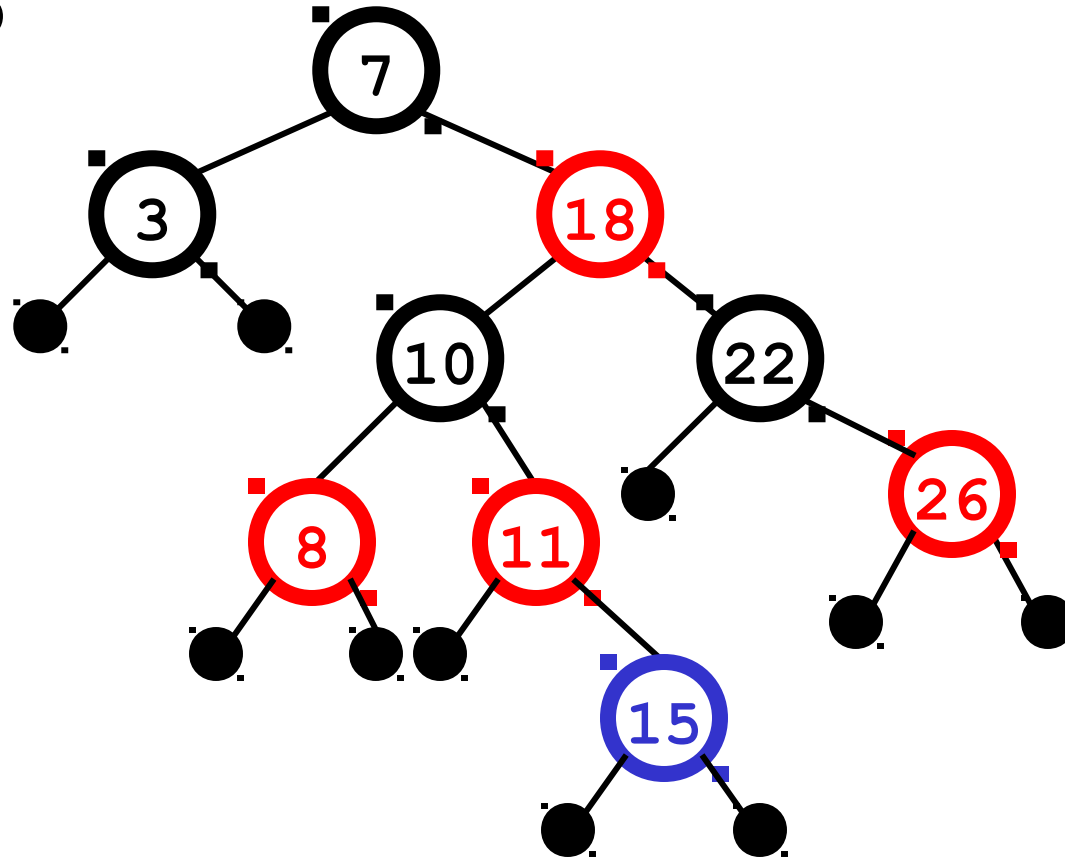
Example

- Insert 10



Example

- Insert 15



Delete BST, p 295 - 298

- BST:
delete x , 3 cases:
 - x has no children
 - x has 1 child
 - x has 2 children
- helper function *transplant*(T, u, v)
 -

	BST- Transplant(T, u, v)
1	if u.p == NIL
2	T.root = v
3	elseif u == u.p.left
4	u.p.left = v
5	else u.p.right = v
6	if v != Nil
7	v.p = u.p

	RB-Transplant(T, u, v)
1	if u.p == T.NIL
2	T.root = v
3	elseif u == u.p.left
4	u.p.left = v
5	else u.p.right = v
6	v.p = u.p

Delete, RB Tree, p 323

- RB-Transplant
- RB-Delete
- RB-DeleteFixup

RB-DELETE(T, z) $y = z$ $y\text{-original-color} = y.\text{color}$ **if** $z.\text{left} == T.\text{nil}$ $x = z.\text{right}$ RB-TRANSPLANT($T, z, z.\text{right}$)**elseif** $z.\text{right} == T.\text{nil}$ $x = z.\text{left}$ RB-TRANSPLANT($T, z, z.\text{left}$)**else** $y = \text{TREE-MINIMUM}(z.\text{right})$ $y\text{-original-color} = y.\text{color}$ $x = y.\text{right}$ **if** $y.p == z$ $x.p = y$ **else** RB-TRANSPLANT($T, y, y.\text{right}$) $y.\text{right} = z.\text{right}$ $y.\text{right}.p = y$ RB-TRANSPLANT(T, z, y) $y.\text{left} = z.\text{left}$ $y.\text{left}.p = y$ $y.\text{color} = z.\text{color}$ **if** $y\text{-original-color} == \text{BLACK}$ RB-DELETE-FIXUP(T, x)

```

while  $x \neq T.root$  and  $x.color == BLACK$ 
    if  $x == x.p.left$ 
         $w = x.p.right$ 
        if  $w.color == RED$ 
             $w.color = BLACK$  // case 1
             $x.p.color = RED$  // case 1
            LEFT-ROTATE( $T, x.p$ ) // case 1
             $w = x.p.right$  // case 1
        if  $w.left.color == BLACK$  and  $w.right.color == BLACK$ 
             $w.color = RED$  // case 2
             $x = x.p$  // case 2
        else if  $w.right.color == BLACK$ 
             $w.left.color = BLACK$  // case 3
             $w.color = RED$  // case 3
            RIGHT-ROTATE( $T, w$ ) // case 3
             $w = x.p.right$  // case 3
             $w.color = x.p.color$  // case 4
             $x.p.color = BLACK$  // case 4
             $w.right.color = BLACK$  // case 4
            LEFT-ROTATE( $T, x.p$ ) // case 4
             $x = T.root$  // case 4
        else (same as then clause with “right” and “left” exchanged)
             $x.color = BLACK$ 

```

Examples

Delete each of these
from the original

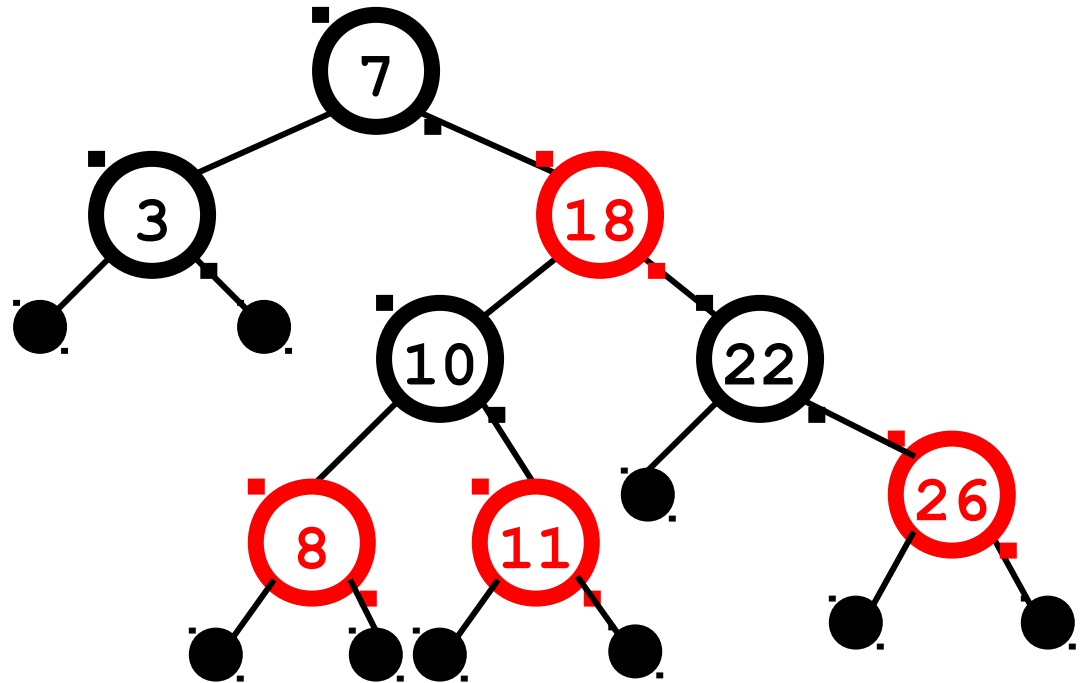
Delete 26

Delete 22

Delete 10

Delete 18

Delete 3



Notes
