
Recurrence Relations – Running Time for Recursive Functions

Chapter 2

Gnome Sort - *trivia*



<http://www.portlandoctopus.com/top-5-garden-gnomes/>

Divide and Conquer Algorithms

- Analysis of divide and conquer algorithms requires knowledge of:
 - Mathematical Induction
 - Substitution/Iterative Method
 - Recurrences

```
class Tree
{


---


public:
    // returns true if t represents a binary
    // search tree containing no duplicate values;
    bool IsBST();

    // return true if & only if all values in the tree are
    // less than val
    bool isLessThan(int val);
    // see above
    bool isGreaterThan(int val);

private:
    int mInfo;
    Tree * mpsLeft;
    Tree * mpsRight;
};
```

Thank you Owen Astrachan

```
// returns true if t represents a binary
// search tree containing no duplicate values;
```

```
bool IsBST()
```

```
{
    bool bLeftIsTree = true, bRightIsTree = true;
    bool bLessThan = true, bGreaterThan = true;
    if( t->left )
    {
        bLeftIsTree = t->left->IsBST();
        bLessThan = t->left->isLessThan(t->info);
    }
    if( t->right )
    {
        bRightIsTree = t->right->IsBST();
        bGreaterThan = t->right->isGreaterThan(t->info);
    }
    return bLessThan &&
           bGreaterThan &&
           bLeftIsTree &&
           bRightIsTree;
} // Complexity with n nodes in the tree?
```

Another Example

- What is the asymptotic complexity of the function below? Assume Combine is $O(n)$

```
// postcondition: a[left] <= ... <= a[right]
void DoStuff(vector<int> & a, int left, int right)
{
    int mid = (left + right)/2;
    if (left < right)
    {
        DoStuff(a, left, mid);
        DoStuff(a, mid + 1, right);
        Combine(a, left, mid, right);
    }
}
```

Recurrence Relation

- A ***recurrence relation*** contains two equations
 - One for the general case
 - One for the base case

Efficiency of Binary Search

Merge Sort

- `MERGE-SORT(A, p, r) // A:Array; p,r: ints`
`// p & r are indices into the array (p < r)`
`if p < r //Check for base case`
`q = $\lfloor (p + r) / 2 \rfloor$ // Divide`
`MERGE-SORT(A, p, q) //Conquer`
`MERGE-SORT(A, q + 1, r) //Conquer`
`MERGE(A, p, q, r) //Combine`

Recurrence Relation

- Let $T(n)$ be the time for Merge-Sort to execute on an n element array.
- The time to execute on a one element array is $O(1)$
- Then we have the following relationship:

Merge Sort

- To solve the recurrence relation we'll write n instead of $O(n)$ as it makes the algebra simpler:
 - $T(n) = 2 T(n/2) + n$
 - $T(1) = 1$
- Solve the recurrence by iteration (substitution)
- Use induction to prove the solution is correct

Recurrence Relations to Remember

$$T(n) = T(n/2) + O(1)$$

$$T(n) = T(n-1) + O(1)$$

$$T(n) = 2 T(n/2) + O(1)$$

$$T(n) = T(n-1) + O(n)$$

$$T(n) = 2 T(n/2) + O(n)$$

Approaches to Algorithm Design

- Incremental
 - Job is partly done – do a little more, repeat until done.
- Divide-and-Conquer (recursive)
 - Divide problem into sub-problems of the same kind.
 - For small subproblems, solve, else, solve them recursively.
 - Combine subproblem solutions to solve the whole thing.

Your Turn

- Solve the following recurrence relation using the expansion (iteration) method
 - $T(n) = T(n-1) + 2n - 1$
 - $T(0) = 0$

For Next Time

- So far we've covered chapters 1, 2, and 3.