# Algorithm Design and Analysis

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Office Hours: MW 11am-noon Friday 1-3pm

CS380 Algorithm Design and Analysis

# Syllabus

- Book
- Schedule
- Grading Assignments/Exams/Quizzes
- Policies
  - Late Policy
  - Grade Complaints
- Moodle

## Overview

- Topics
- Goals
- Where does this class fit?
  - Data Structures  $\rightarrow$  Algo  $\leftarrow$  Discrete Math
    - Big Oh, Data Structs
      Graphs, Proofs, Logic
  - Theory  $\leftrightarrow$  Algo

### What is an Algorithm? chapter 1

#### How do we discuss algorithms?

#### How to we evaluate algorithms?

• Performance

• and ...

## Why Study Algorithms?

#### Correctness

## Demonstrating Incorrectness

- Searching for counterexamples is the best way to disprove the correctness of a heuristic.
- Think about all small examples.
- Think about examples with ties on your decision criteria (e.g. pick the nearest point).
- Think about examples with extremes of big and small.

## Induction and Recursion

- Failure to find a counterexample to a given algorithm does not mean "it is obvious" that the algorithm is correct.
- Mathematical induction is a very useful method for proving the correctness of recursive algorithms.
- Recursion and induction are the same basic idea: (1) basis case, (2) general assumption, (3) general case.

### Correctness is Not Obvious!

 Visit each point once, minimizing the distance moved



#### Nearest Neighbor Tour



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#### Nearest Neighbor Tour



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## A Correct Algorithm

## Why Not Use a Supercomputer

- A faster algorithm running on a slower computer will always win for sufficiently large instances
- Usually, problems don't have to get that large before the faster algorithm wins

## **Expressing Algorithms**

- What are the possible ways to express an algorithm?
  - English
  - Pseudocode
  - Programming Language

## The RAM Model, section 2.2

- Algorithms can be studied in a machine and language independent way.
- Each "simple" operation (+, -, =, if, call) takes exactly one step.
- Loops and subroutines are not simple operations.
- Each memory access takes one step.

# Tuning

 But what about registers, cache, RAM, Virtual Memory, pipelining, branch prediction, etc.

http://www.dyninst.org/harmony

### Best, Worst, and Average-Case

• Worst case:

Best case:

• Average case \*:

#### rate of growth OR order of growth

## Example: Sorting

- Input: A sequence of n numbers
  <a<sub>1</sub>, a<sub>2</sub>, ..., a<sub>n</sub>>
- Output: A permutation (reordering)
  <a '₁, a '₂, ..., a 'ₙ> of the input sequence
  such that a '₁ ≤ a '₂ ≤ ... ≤ a 'ₙ

 We seek algorithms that are correct and efficient

## Insertion Sort, p 26

- INSERTION-SORT(A) // A is an array
- 1 for j = 2 to A.length
- $2 \quad key = A[j]$
- 3 // Insert A[j] in to the correct location
- 4 i = j 1
- 5 while i > 0 and A[i] > key
- A[i+1] = A[i]
- 7 i = i 1
- 8 A[i+1] = key

## Example

How would insertion sort work on the following numbers?

• 3 1 7 4 8 2 6

### **Insertion Sort**

- Is the algorithm correct?
- How efficient is the algorithm?
- How does insertion sort do on sorted permutations?
- How about unsorted permutations?

## Analysis of Insertion Sort

Best Case

## Analysis of Insertion Sort

Worst Case

- We generally care about this analysis.
  - gives the upper bound time
  - average case is often closer to worst than best
    - and average case is often very hard to compute
  - unless we know we are in a best case

# Your Turn

- Problem: How would insertion sort work on the following characters to sort them alphabetically (from A -> Z)? Show each step.
  - S O R T E D

this is a good exam question.

### For Next Time

Read Chapters 1 and 2 from the book.