

CS310

P vs NP

the steel cage death match

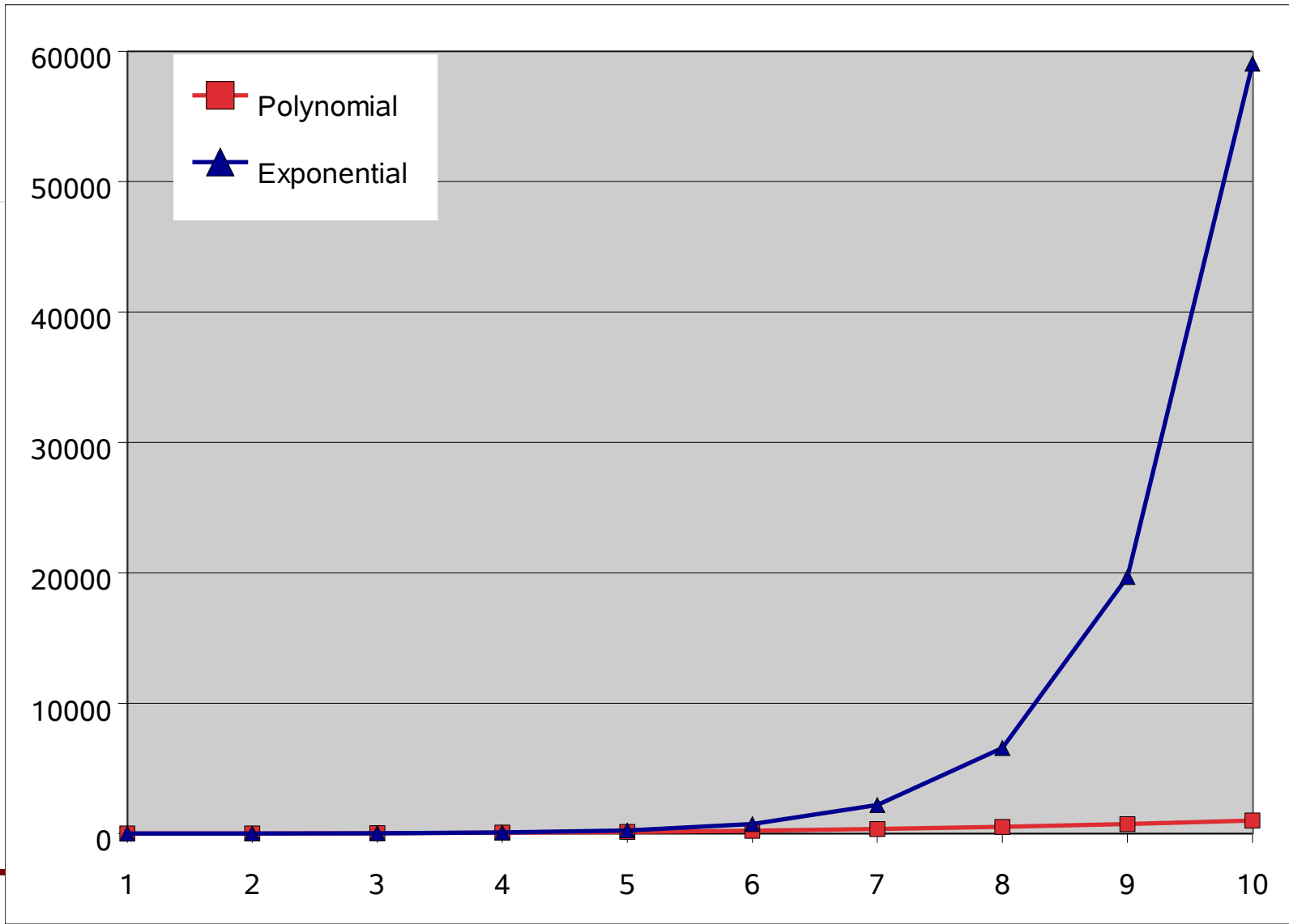
How hard is a problem to solve?

Section 7.2

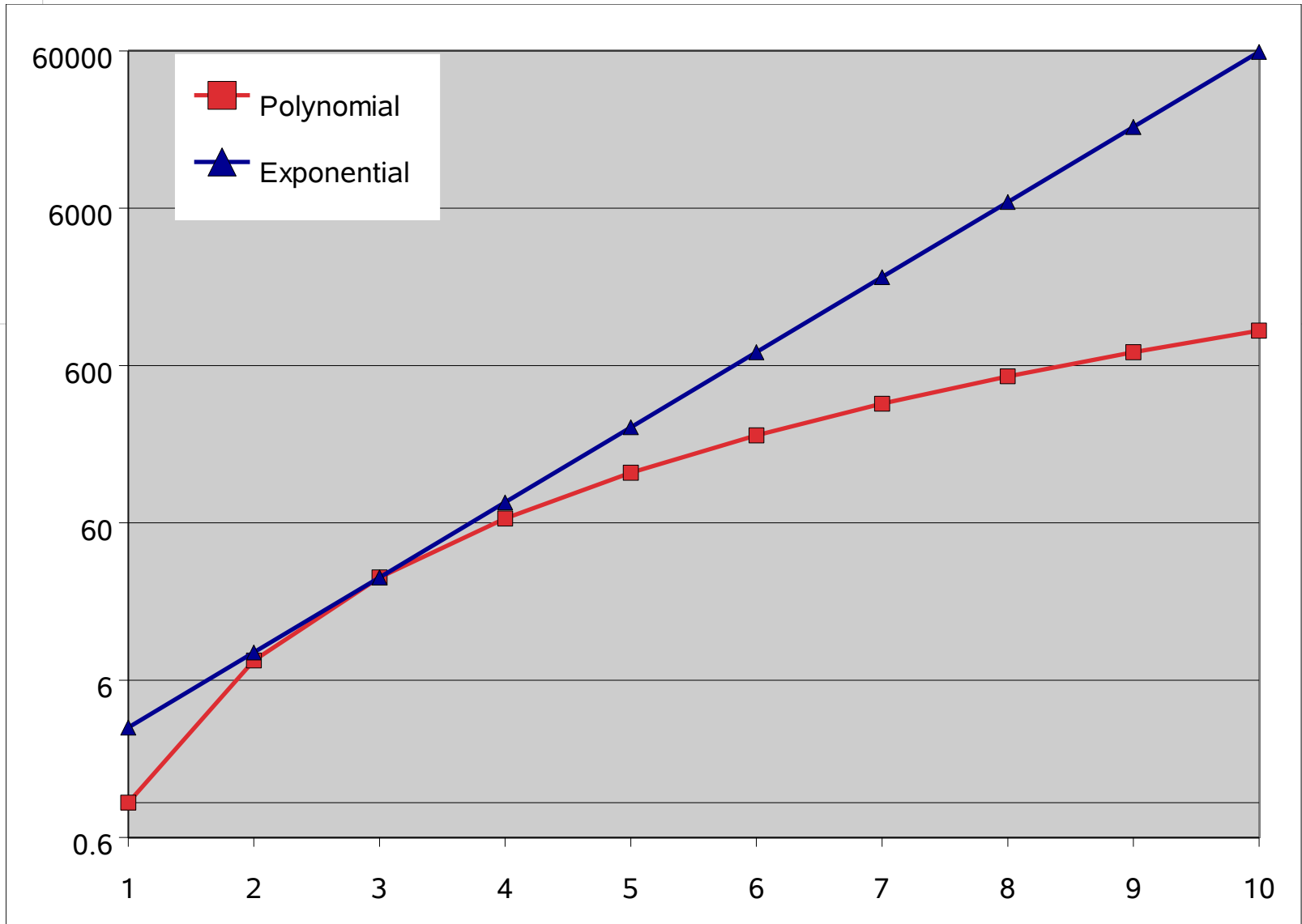
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Polynomial vs Exponential

- Polynomial: n^3 Exponential: 3^n



Log Scale



Complexity relationships between models

- Theorem 7.8: let $t(n) \geq n$, every $t(n)$ time multi-tape TM has an equivalent $O((t(n))^2)$ time single-tape TM.
 - polynomial difference
- Theorem 7.9: Every $t(n) \geq n$ time ND single tape TM has an equivalent $2^{O(t(n))}$ time deterministic single tape TM
 - exponential difference

The class P

- P is the class of languages that are decidable in polynomial time on a **deterministic, single tape** TM
- Problems in class P
 - PATH: { $\langle G, s, d \rangle$ | G is a directed graph, find a directed path from s to d }
 - RELPRIME: { $\langle x, y \rangle$ | x and y are relatively prime }
 - Euclidean algorithm
- Every context-free language is in P

Real Life

- Problems in class P are usually manageable on a real computer
 - n^k
 - though $k=100$ may introduce some practical problems

The class NP

- NP is the class of languages that are decidable in **polynomial time on a nondeterministic single tape TM**
 - Problems in class NP
 - HAMPATH: { $\langle G, s, t \rangle$ | G is a directed graph, with Hamilton path from s to t } (a path that passes through every vertex of a graph exactly once)
 - These problems are decidable on a deterministic single tape TM in exponential time

Verifier

- A verifier of a language, A , is an algorithm, V , such that

$A = \{ w \mid V \text{ accepts } \langle w, c \rangle \text{ for some string } c \}$
where c is a certificate

$|c|$ is polynomial in terms of $|w|$

- NP is the class of languages that have polynomial time (in terms of the length of w) verifiers

P vs NP

- $P \subseteq NP$
 - unknown if the classes are unequal
- If $P = NP$, then all problems in NP can be solved in polynomial time, **if** we are clever enough to find the right algorithm
- NP-Completeness
 - set of problems in NP whose complexity is related to all problems in NP
 - if an NP-Complete problem can be shown to be in P, then $P=NP$
 - boolean satisfiability, for example