CS310

Pumping Lemma

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Non-Regular Languages

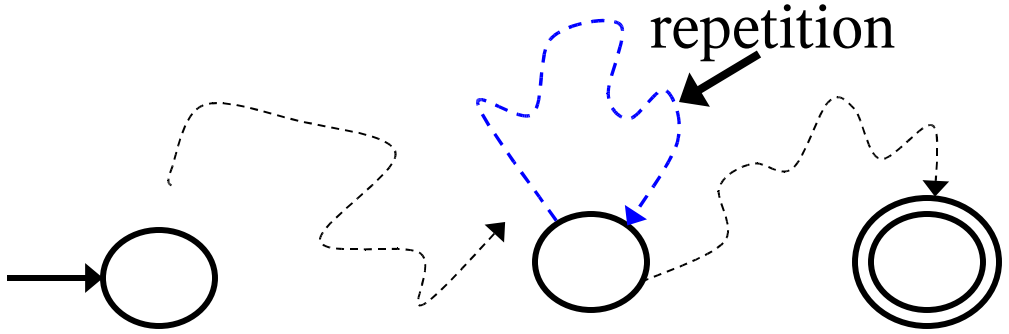
- Languages that *cannot* be represented by a finite automaton
  - Such as?
- How do we prove a language is not regular?

C = { w | w has an equal number of 0s and 1s}

D = { w | w has an equal number of occurrences of 01 and 10 as substrings }

Pumping Lemma (Informal)

Pumping: The length of the string could be 'pumped' up by repeating the cycle, *and the string would still be accepted.*



- All regular languages have a property
  - *the pumping length, p*
- |w| = n, how many states do we go through?

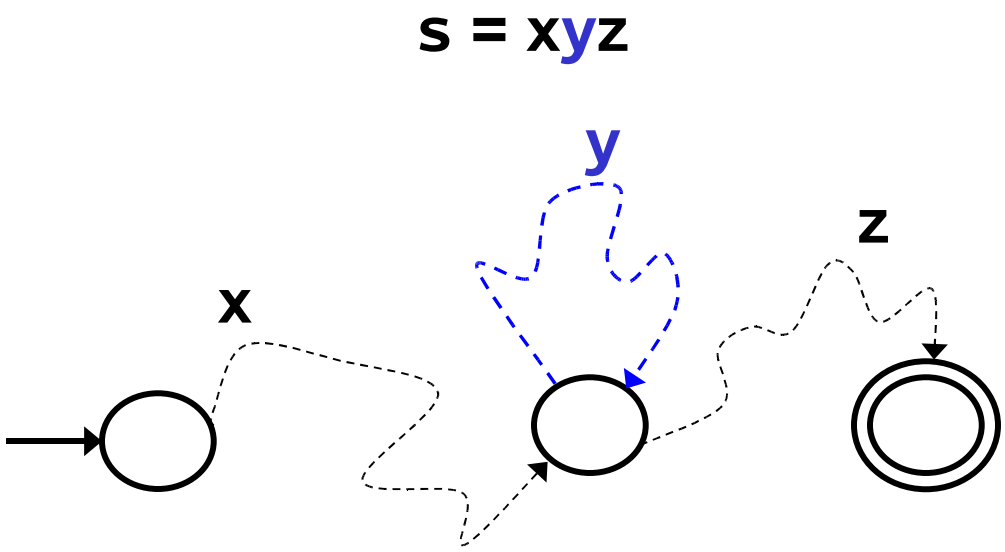
Pumping Lemma (Formally)

- DFA: $M = (Q, \Sigma, \delta , q_0, F)$

If $|Q| = p$ and $s \in L(M)$ and $|s| >= p$

then there exists at least one state that was

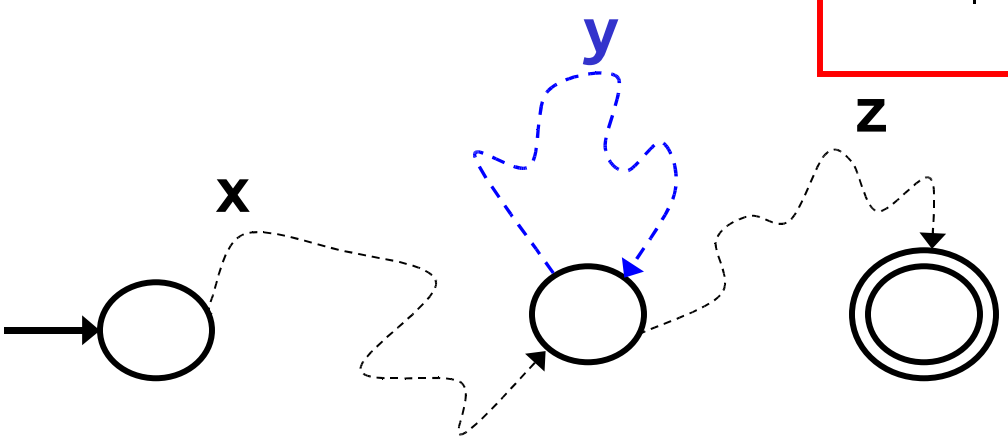
visited twice within the first *p* input symbols



Pumping Lemma (Formally)

- If A is a regular language, then :

- $i >= 0$, $xy^iz \in L(M)$
- $|y| > 0$ (x, z may be ε)
- $|xy| <= p$



Pumping Lemma In Action

- Find a string, s ϵ L, |s| >= p, that cannot be pumped to show language L is not regular.
  - Find a string that exhibits the "essence" of nonregularity
  - Proof method?

- L = { w | w contains equal number of 0s and 1s }

Pumping Lemma in Action

- L = { w | w contains equal number of 0s and 1s }

  use a different string:

  Can that be pumped?

s =

x =

y =

z =

Practice

- $L = \{ ww \mid w \in \{0, 1\}^* \}$

  What string should we chose?

  what does ww mean?

  Can that be pumped?