Minimum Spanning Trees

Chapter 23

Spanning Tree

- What are the edges you need to keep the graph connected?
  - If you remove any edge, the graph becomes disconnected
- Minimum Spanning Tree
  - minimize the total weight of the edges

Problem

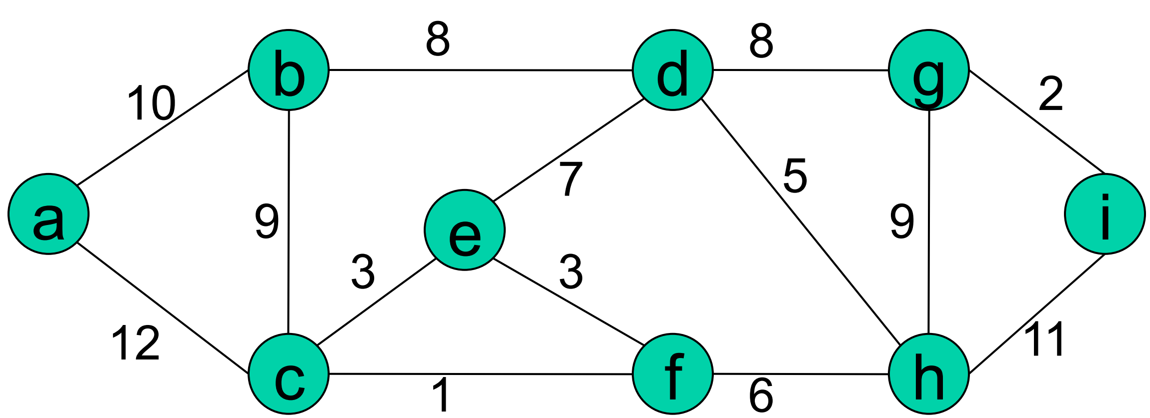
- A town has a set of houses and a set of roads
- A road connects 2 and only 2 houses
- A road connecting houses u and v has a repair cost w(u, v)
- Goal: Repair enough (and no more) roads such that
  - Everyone stays connected
  - Total repair cost is minimum

Minimum Spanning Tree

- Model as a graph:
  - Undirected graph G = (V, E)
  - Weight w(u, v) on each edge (u, v) in E
  - Find T that is a subset of E such that
    - T connects all vertices, and
    - $w(T) = \sum_{(u,v)\in T} w(u,v)$ is minimized

Minimum Spanning Tree

- A spanning tree whose weight is minimum over all spanning trees is called a minimum spanning tree
- Example



Growing an MST

- Properties of an MST

- Building up a Solution

Generic MST Algorithm

GENERIC-MST(G, w)

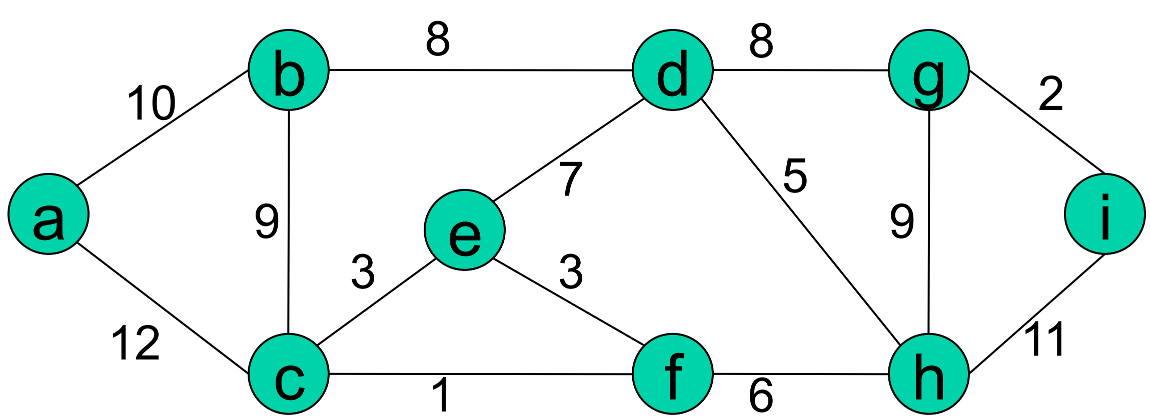
$A = \emptyset$
while A is not a spanning tree
    find an edge (u, v) that is safe for A
    $A = A \cup \{(u, v)\}$
return A

Proof via Loop Invariant

- Initialization
- Maintenance
- Termination

Finding a Safe Edge

- How do we find safe edges?
- Looking at the example below, Edge (c,f) has the lowest weight of any edge in the graph. Is it safe for A?



Finding a Safe Edge

- Intuitively: Let S, a subset of V, be any set of vertices that includes c but not f (f is in V-S).
- In any MST, there has to be one edge that connects S with V-S.
- Why not choose the edge with the minimum weight?

Definitions

- Let S be a subset of V and A be a subset of E
  - A **cut** (S, V-S) is a partition of vertices into disjoint sets V and S-V
  - Edge (u,v) in E **crosses** cut (S,V-S) if one endpoint is in S and the other is in V-S
  - A cut **respects** A if and only if no edge in A crosses the cut
  - An edge is a **light edge** crossing a cut if and only if its weight is minimum over all edges crossing the cut

Theorem

- Let A be a subset of some MST, (S,V-S) be a cut that respects A, and (u,v) be a light edge crossing (S,V-S).
- Then….

Generic-MST

- So, in a generic MST
  - A is a forest containing connected components. Initially, each component is a single vertex
  - Any safe edge merges two of these components into one. Each component is a tree
  - Since an MST has exactly |V|-1 edges, the for loop iterates |V|-1 times. Equivalently, after adding |V|-1 safe edges, we're down to just one component

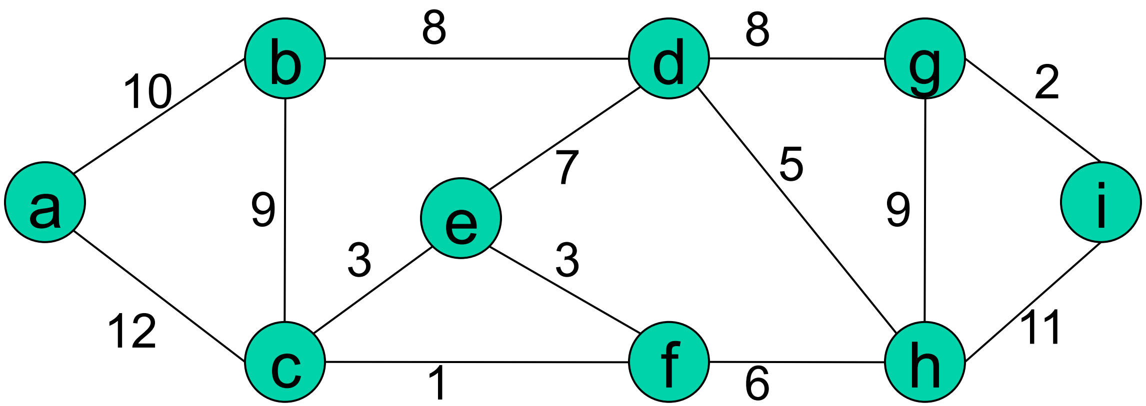
Kruskal’s Algorithm

- G = (V,E) is a connected, undirected, weighted graph. w:E->R
  - Starts with each vertex being its own component
  - Repeatedly merges two components into one by choosing the light edge that connects them
  - Scans the set of edges in monotonically increasing order by weight
  - Uses a disjoint-set data structure to determine whether an edge connects vertices in different components

Kruskal(V,E,w)

```
KRUSKAL(G, w)
A = ∅
for each vertex v ∈ G.V
    MAKE-SET(v)
sort the edges of G.E into nondecreasing order by weight w
for each (u, v) taken from the sorted list
    if FIND-SET(u) ≠ FIND-SET(v)
        A = A ∪ {(u, v)}
        UNION(u, v)
return A
```

Example



Prim’s Algorithm

- Builds one tree, so A is always a tree
- Starts from an arbitrary “root” r
- At each step, find a light edge crossing cut $(V_A, V-V_A)$, where V_A = vertices that A is incident on. Add this edge to A

How to Find a Light Edge Quickly

- Use a priority queue Q:
  - Each object is a vertex in $V\text{-}V_A$
  - Key of v is minimum weight of any edge (u,v), where u is in V_A
  - Then the vertex returned by EXTRACT-MIN is v such that there exists u in V_A and (u,v) is a light edge crossing (V_A , $V\text{-}V_A$)
  - Key of v is infinity if v is not adjacent to any vertices in V_A

Prim's Algorithm

- The edges of A will form a rooted tree with root r:
  - r is given as an input to the algorithm, but it can be any vertex
  - Each vertex knows its parent in the tree by the attribute $\pi[v]$ = parent of v. $\pi[v]$ = NIL if v = r or v has no parent

PRIM(G,w,r)

$\textsc{Prim}(G, w, r)$

$Q = \emptyset$

for each $u \in G.V$

    $u.key = \infty$

    $u.\pi = \text{NIL}$

    $\textsc{Insert}(Q, u)$

$\textsc{Decrease-Key}(Q, r, 0)$ // $r.key = 0$

while $Q \neq \emptyset$

    $u = \textsc{Extract-Min}(Q)$

    for each $v \in G.Adj[u]$

        if $v \in Q$ and $w(u, v) < v.key$

            $v.\pi = u$

            $\textsc{Decrease-Key}(Q, v, w(u, v))$

Example