
Elementary Graph Algorithms

Chapter 22

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Topological Sort

- A topological sort is performed on a directed acyclic graph
- A topological sort is a linear ordering of all vertices of a graph such that if G contains an edge (u, v) , then u appears before v in the ordering

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Topological Sort

- A topological sort of a graph can be viewed as an ordering of its vertices along a horizontal line so that all directed edges go from left to right
- Directed Acyclic Graphs (DAG) are used in many applications to indicate precedences among events
- What is a DAG?

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Topological Sort

- Good for modeling processes and structures that have a partial order:
 - $a > b$ and $b > c$ implies that $a > c$
 - But may have a and b such that neither $a > b$ nor $b > a$

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TOPOLOGICAL-SORT(G)

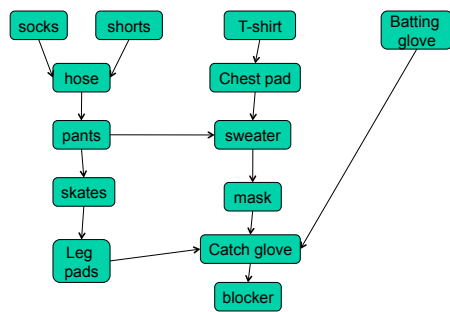
- Call DFS(G) to compute finishing times $f[v]$ for each vertex v
- As each vertex is finished, insert it onto the front of a linked list
- Return the linked list of vertices

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Example



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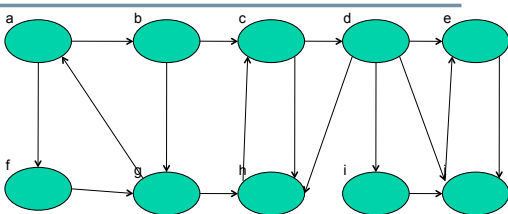
Topological Sort

- Running time for topological sort is:

Strongly Connected Components

- Given a directed graph $G = (V, E)$
- A strongly connected component (SCC) of G is a maximal set of vertices $C \subseteq V$
- Such that for all $u, v \in C$ both $u \rightarrow v$ and $v \rightarrow u$

Example



- Identify the strongly connected components

Transpose

- Algorithm uses $G^T = \text{transpose of } G$
 - G^T
- How long does it take to create G^T if using adjacency lists?
- Observation: G and G^T have the same SCC's.

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SCC(G)

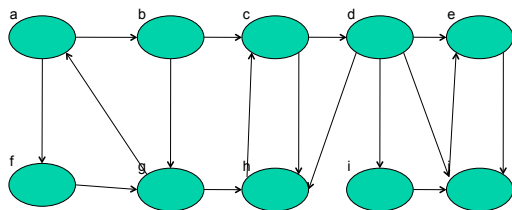
- Call $\text{DFS}(G)$ to compute finishing times $f[u]$ for all u
- Compute G^T
- Call $\text{DFS}(G^T)$, but in the main loop, consider vertices in order of decreasing $f[u]$ (as computed in first DFS)
- Output the vertices in each tree of the depth-first forest formed in second DFS as a separate SCC

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Example



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