

---

## Elementary Graph Algorithms

### Chapter 22

---

---

---

---

---

---

---

---

---

## Graph Representation

- Given a graph  $G = (V, E)$
- The graph may be directed or undirected
- There are two common ways to represent for algorithms:
  - Adjacency lists.
  - Adjacency matrix.

---

---

---

---

---

---

---

---

---

## Running Times

- We will be talking about the running time of graph algorithms in terms of both Vertices  $|V|$  and Edges  $|E|$
- We can remove the cardinality when in asymptotic notation
  - Example:  $O(V + E)$

---

---

---

---

---

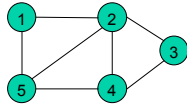
---

---

---

## Adjacency Lists

- Array Adj of  $|V|$  lists, one per vertex
- Vertex  $u$ 's list has all vertices  $v$  such that  $(u, v) \in E$
- Example:



4/1/09

CS380 Algorithm Design and Analysis

4

---

---

---

---

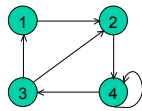
---

---

---

---

## Example



- Space:
- Time to list all vertices adjacent to  $u$ :
- Time to determine if  $(u, v)$  is an edge:

4/1/09

CS380 Algorithm Design and Analysis

5

---

---

---

---

---

---

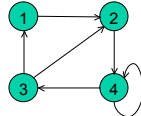
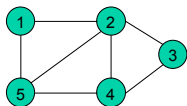
---

---

## Adjacency Matrix

- $|V| \times |V|$  matrix  $A = (a_{ij})$

$$a_{ij} = \begin{cases} 1 & \text{if } (i, j) \in E \\ 0 & \text{otherwise} \end{cases}$$



4/1/09

CS380 Algorithm Design and Analysis

6

---

---

---

---

---

---

---

---

## Adjacency Matrix

---

- Space:
- Time to list all vertices adjacent to  $u$ :
- Time to determine if  $(u,v)$  is an edge:
  
- What about weighted graphs?

4/1/09

CS380 Algorithm Design and Analysis

7

---

---

---

---

---

---

---

---

## Breadth-First Search

---

- Input: Graph  $G = (V, E)$ , either directed or undirected, and source vertex  $s$  is in  $V$ .
- Output:
  - $d[v]$  = distance (smallest # of edges) from  $s$  to  $v$ , for all  $v$  in  $V$ .
  - $\pi[v] = u$  such that  $(u,v)$  is last edge on shortest path  $s \rightarrow v$
- $u$  is  $v$ 's predecessor
- Set of edges  $\{(\pi[v],v) : v \neq s\}$  forms a tree

4/1/09

CS380 Algorithm Design and Analysis

8

---

---

---

---

---

---

---

---

## Breadth-First Search

---

- Idea: Send a wave out from  $s$ .
  - First hits all vertices 1 edge from  $s$ .
  - From there, hits all vertices 2 edges from  $s$ .
  - Etc.
- Use FIFO queue  $Q$  to maintain wavefront.
  - $v$  is in  $Q$  if and only if wave has hit  $v$  but has not come out of  $v$  yet

4/1/09

CS380 Algorithm Design and Analysis

9

---

---

---

---

---

---

---

---

## BFS( $G, s$ )

---

---

---

---

---

---

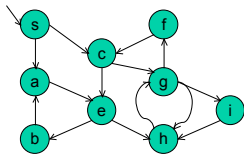
---

---

---

## Example

---



---

---

---

---

---

---

---

---

## Breadth-First Search

---

- Will breadth-first search reach all vertices?
  
- Time =  $O(\quad)$

---

---

---

---

---

---

---

---

## Depth-First Search

---

- Input:  $G = (V, E)$ , directed or undirected. No source vertex given.
- Output: 2 timestamps on each vertex:
  - $d[v]$  = discovery time
  - $f[v]$  = finishing time
  - $\pi[v] = u$  such that  $(u,v)$  is last edge on shortest path  $s \rightarrow v$

---

---

---

---

---

---

---

---

## DFS(G)

---

---

---

---

---

---

---

---

---

## DFS-VISIT(u)

---

---

---

---

---

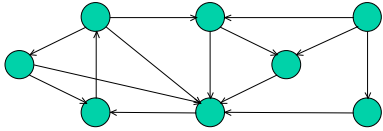
---

---

---

---

## Example



---

---

---

---

---

---

---

---

## Depth-First Search

- Running Time =

---

---

---

---

---

---

---

---

## Classification of Edges

- Tree edge:
- Back edge:
- Forward edge:
- Cross edge:

---

---

---

---

---

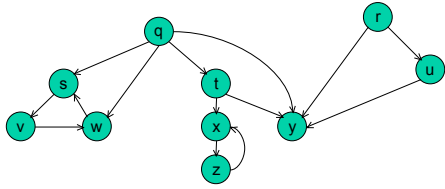
---

---

---

## Your Turn

- Solve exercise 22.3-2 on page 547



---

---

---

---

---

---

---

---