

CS430 Computer Architecture

Spring 2015

CS430 - Computer Architecture

Chapter 12 Types of Operations

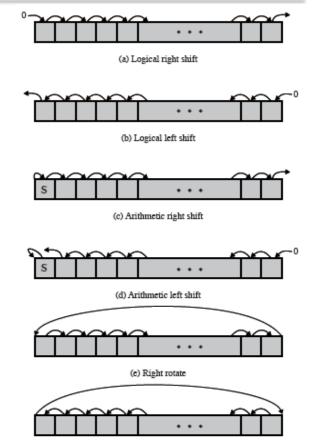
- Data Transfer copies data from a source operand into a destination operand
- x86 Examples (Reference: <u>http://zeus.cs.pacificu.edu/ryand/cs320/2005/cs320.html</u>)
 - > mov ax,1 ; move 1 into ax
 - > movzx ax, 1000000b ; mov 128 zero-extended into ax
 - > movsx ax, 1000000b ; mov -128 sign-extended into ax
 - > push ebx ; push 32-bit contents of ebx onto the stack
 - > pop edx

- Arithmetic perform some arithmetic calculation and in the case where the processor has a flags register, sets the flags accordingly
- x86 Examples
 - > add ax,bx ; ax<-ax+bx</pre>
 - > sub ax,1 ; ax<-ax-1</pre>
 - inc cx

> dec cx

- Logical instructions used to do some kind of bit manipulation
- x86 Examples
 - > and bh,0fh
 - > or ax,10h
 - > xor ax,bx

- More Logical
 - > shr ax,1 ; (a)
 - > shl ax,1 ; (b)
 - > sar bh,1 ; (c)
 - > sal bh,1 ; (d)
 - > ror edx,1 ; (e)
 - > rol edx,1 ; (f)



(f) Left rotate

• Transfer of Control

- Conditional branch (conditional jump)
- > Unconditional branch
- Subroutine call
- > Interrupt

- Conditional branch branching is conditionally based on some flags register or some status register
- x86 Example
 - > jne top ; branch to top if ZF = 0
 - jb top ; unsigned ... branch to top if not above or equal ; CF = 1
 - jl top ; signed ... branch to top if not greater or equal ; SF <> OF

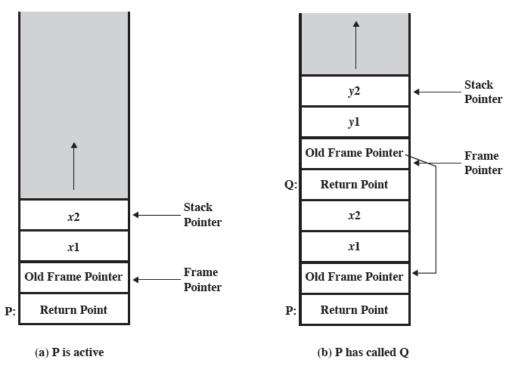
- Conditional branch
- x86 Example
 - Conditional branch instructions assume a calculation occurred setting flags in the flags register BEFORE the branch occurs
 - dec ax jne top

- Unconditional branch the branch occurs regardless
- x86 Example
 - jmp top

- Subroutine call
- x86 Example

> call Foo ; Foo is an assembly language subroutine

Subroutine call - what is happening below?



CS430 - Computer Architecture

 Subroutine call - a typical x86 procedure might begin with the following code:

push ebp

mov ebp, esp

sub esp, space_for_locals

- Interrupt
- x86 Example
 - int Oh ; transfer control to the address stored in the ; interrupt vector table at location 0

- Interrupt in x86 Real Mode,
 - an interrupt has an integer in the range of 0-255 called the interrupt type
 - The addresses from 00000 to 003ff are reserved for interrupt vectors
 - An interrupt vector is an address (segment & offset) of a particular interrupt service routine

• Interrupt Vector Type

 Interrupt Vector Type
 Stored At

 0
 00000:00003

 1
 00004:00007

 4t : 4t+3

 We see that each interrupt vector requires four bytes. The first two bytes contain the offset (bytes reversed). The next two bytes contain the segment (bytes reversed).

- How do interrupts work? Remember, after most instructions, the microprocessor checks for pending interrupts. If detected, the microprocessor
 - 1) Push the flags register on the stack (why?)
 - 2) Clear the interrupt and trap flags (why?)
 - 3) Push CS
 - 4) Determine the interrupt location based on the type
 - 5) CS = second word of the interrupt vector
 - 6) Push IP
 - 7) IP = first word of the interrupt vector

• The instruction IRET transfers control back to the caller