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# CS310

## Converting NFA to DFA

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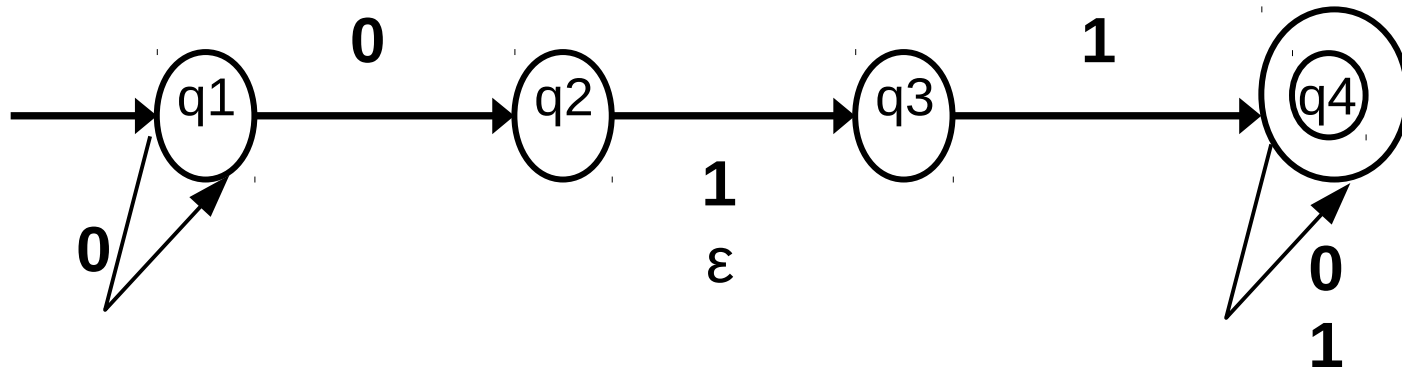
September 10, 2014

# Quick Review

- 5 tuple (  $Q, \Sigma, \delta, q_0, F$  )

$$\Sigma_\epsilon = \Sigma \cup \{\epsilon\}$$

$$\delta : Q \times \Sigma_\epsilon \rightarrow P(Q)$$



# Convert NFA to DFA

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- Two machines are equivalent if they recognize the same language
- Every NFA has an equivalent DFA (Th 1.39)  
$$\delta_{nfa} : Q \times \Sigma_{\varepsilon} \rightarrow P(Q)$$
- The DFA will need to represent all subsets in  $P(Q)$  (how many?)
  - let's assume no  $\varepsilon$ -transitions initially

# Convert NFA to DFA

- NFA is  $N = (Q, \Sigma, \delta, q_0, F)$
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- DFA is  $M = (Q', \Sigma', \delta', q_0', F')$

$Q' =$

$q_0' =$

$F' =$

$\delta':$

NFA

# Example (without $\epsilon$ or $\delta_{dfa}$ )

DFA

$Q = \{q_0, q_1\}$

$\Sigma = \{a, b\}$

$Q_0 = q_0$

$F = \{q_0\}$

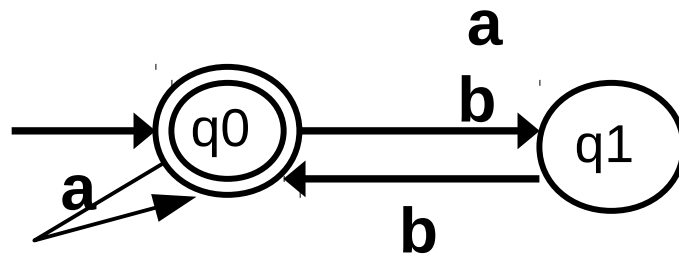
$Q = \{q_0, q_1\}$

$\Sigma = \{a, b\}$

$Q_0 = q_0$

$F = \{q_0\}$

$\delta$	a	b
q0	{q0, q1}	{q1}
q1	{}	{q0}



Let's define the  $\delta_{dfa}$  (still no  $\varepsilon$ )

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$\delta_{nfa} : Q \times \Sigma_\varepsilon \rightarrow P(Q)$  in NFA

$\delta_{dfa} : Q' \times \Sigma \rightarrow Q'$  in DFA

$R \in Q', a \in \Sigma$

$\delta_{dfa}(R, a) =$

# Converting NFA to DFA - $\epsilon$ Transitions

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- Define start state and  $\delta_{dfa}$  to include all states that can be reached from a given state by 0 or more  $\epsilon$  transitions

# Conversion Example (with $\epsilon$ )

DFA

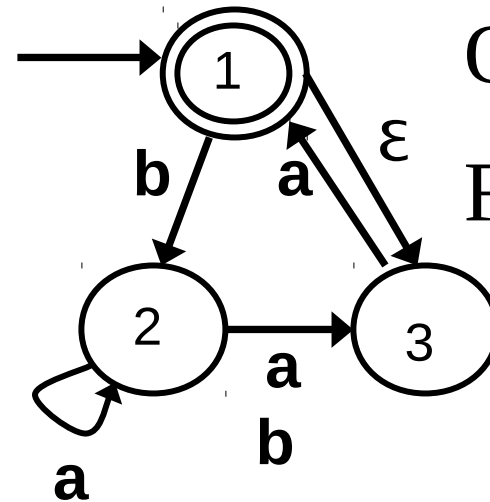
$Q = \{1, 2, 3\}$

$\Sigma = \{a, b\}$

$Q_0 = 1$

$F = \{1\}$

$\delta_{dfa} =$



$Q = \{1, 2, 3\}$

$\Sigma = \{a, b\}$

$Q_0 = 1$

$F = \{1\}$