CS310

Finite Automata

Sections:1.1, 1.2 page 44

September 5, 2014

Quick Review

- Deterministic Finite Automata:

  5-tuple $(Q, \sum, \delta, q_0, F)$

  Q: finite set of states

  $\sum$: alphabet (finite set)

  δ : transition function (δ: Qx$\sum$ **->** Q)

  q_0: start state

  F : accepting states (subset of Q)

- Language A is ***regular*** if there exists a Finite Automata that recognizes A.

Regular Language

- Determinism?

- Regular language
  - Example?

  - Example of non-regular language?

Regular Operations on Languages

- Given two languages, A,B, we can create new *languages* in a variety of ways:
  - What operations have we seen?

Examples

$\Sigma=\{0,1\}$ A=$\{w|$ w ends in 1$\}$
B=$\{w \mid$ w begins with 00$\}$

$A \cup B$ =

AB =

A^* =

$A \cap B$ =

$\bar{A}$ =

Closure of Regular Languages

- A set is ***closed*** under some operation, Examples?

- ***Regular operations***

Proof

- Theorem 1.25: The class of regular languages is closed under the union operation.

  If A and B are regular languages, so is $A \cup B$

  What do we need to prove?

  What does regular mean?

  What does it mean for $A \cup B$ to be regular?

Build the machine

$\sum=\{0,1\}$

$A=\{w|$ w contains a 1 in the penultimate position$\}$

A = { }

Nondeterminism

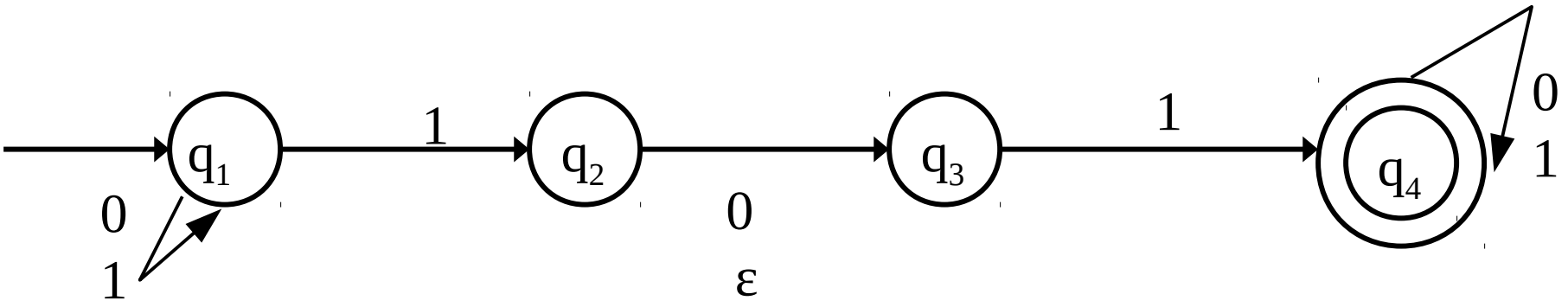
- Nondeterministic Finite Automata:

NFA

- ε transitions

- Why would we ever use this?

Example



- Does this NFA accept 010110?
- What sequence of states does it go through?

Build the machine

$\sum=\{0,1\}$

$A=\{w|$ w contains a 1 in the penultimate position$\}$

A = { }

Proof

- Theorem 1.26: The class of regular languages is closed under the concatenation operation.

  If A and B are regular languages, so is AB.

  What do we need to prove?

  What does regular mean?

  What does it mean for AB to be regular?

  Problems?

Examples

A = {north, northern} B= {east, west}

s = northerneast is in AB

Are A and B regular languages?

A = {w | w = begins with 1 ends with 0}

B = {w | w = begins with 0 ends with 1}

s = 1000011