

CS310

NP-Completeness

Section 7.4

November 30, 2014

NP-Complete

- NP-Completeness
 - set of problems in NP whose complexity is related to all problems in NP
 - if an NP-Complete problem can be shown to be in P, then $P=NP$
 - boolean satisfiability, for example
 - vertex-cover
 - clique
 - Hamilton Path

Boolean Satisfiability

And
Or
Not

Like a circuit

- Is a boolean formula satisfiable?
 - Does some set of values produce true?

$$\phi = (\bar{x} \vee y \vee z) \wedge (x \vee \bar{z} \vee y) \quad \bar{z} \text{ means } \neg z$$

SAT = { $\langle \phi \rangle$ | ϕ is a satisfiable Boolean formula }

- Clause
- Conjunctive normal form (cnf)
- 3cnf

3SAT = { $\langle \phi \rangle$ | ϕ is a satisfiable 3cnf Boolean formula }

- Cook-Levin Theorem: SAT \in P iff P = NP

Reducibility

- If problem A is *efficiently* reducible to problem B , an efficient solution to B can be used to solve A *efficiently*

A function f is a *polynomial time computable function* if some polynomial time TM exists that halts with just $f(w)$ on the tape when run on input w .

Cont.

- Language A is polynomial time reducible to language B , $A \leq_p B$, if a polynomial time computable function f exists where for every w :

$$w \in A \Leftrightarrow f(w) \in B$$

3SAT reduces to CLIQUE

- Polynomial time reduction
- If CLIQUE is in P, so is 3SAT

- Turn 3SAT into a graph
 - Identify a CLIQUE to find a solution to 3SAT

NP-Complete

- B is NP-Complete if:
 - B is in NP
 - Every A in NP is polynomial time reducible to B
- If B is NP-Complete, and $B \in P$, then $P=NP$

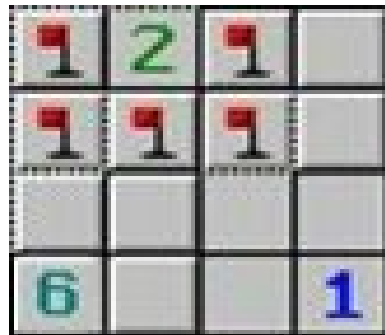
SAT is NP-Complete

- SAT is in NP
- Show that every language in NP can be polynomial time reduced to SAT

Minesweeper is NP-Complete

- Given a partial board, is it a valid Minesweeper board?

Can convert SAT problem into a Minesweeper board.



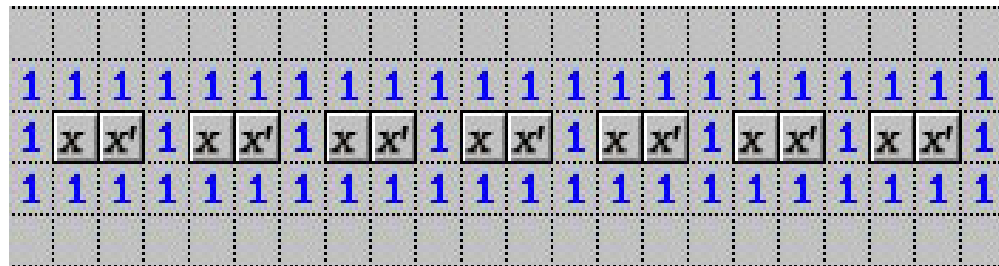
Invalid Board



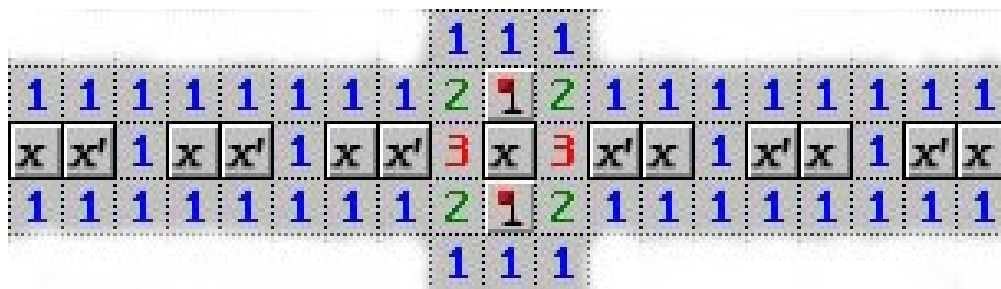
http://www.claymath.org/Popular_Lectures/Minesweeper/

Build the board from SAT

- Cell with mine is True



Wire to propagate a value



Not Gate

http://www.claymath.org/Popular_Lectures/Minesweeper/

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Cont.

- And Gate

U & V are input wires, W is output Wire

| | | | | | | | | | | | | | | | | | | | | | | | |
|--------|---|----|---|---|----|----|---|----|----------------|----------------|----------------|----|---|----|----|---|----|---|---|----|-----|-----|-----|
| U ↓ | ⋮ | ⋮ | ⋮ | | | | | | | | | | | | | | | | | | | | |
| | 1 | 1 | 1 | | | 1 | 2 | 2 | 1 | | 1 | 1 | 1 | | 1 | 1 | 1 | | | | | | |
| | 1 | u' | 1 | | | 2 | 1 | 1 | 3 | 2 | 3 | 1 | 2 | 1 | 2 | 1 | 3 | 2 | 1 | | | | |
| | 1 | u | 1 | 1 | 2 | 4 | 1 | s | a ₁ | a ₂ | a ₃ | t' | 3 | t | t' | 3 | 1 | 1 | 2 | | | | |
| 1 | 2 | 2 | 1 | 1 | 1 | 1 | 4 | 1 | 3 | 2 | 3 | 1 | 2 | 1 | 1 | 2 | t | 1 | 2 | | | | |
| 2 | 1 | u' | 2 | 2 | 4 | s' | 3 | 1 | 1 | 0 | 1 | 1 | 1 | 0 | 0 | 1 | 2 | 2 | 1 | | W → | | |
| 2 | 1 | 1 | 3 | u | u' | s | 2 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | t' | 1 | 1 | 1 | 1 | ... | |
| 2 | 4 | 5 | 1 | 4 | 1 | 4 | t | t' | 1 | t | t' | 1 | t | t' | 1 | t | 2 | t | 1 | t' | t | 1 | ... |
| 2 | 1 | 1 | 3 | v | v' | r | 2 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | t' | 1 | 1 | 1 | 1 | 1 | ... |
| 2 | 1 | v' | 2 | 2 | 4 | r' | 3 | 1 | 1 | 0 | 1 | 1 | 1 | 0 | 0 | 1 | 2 | 2 | 1 | | | | |
| 1 | 2 | 2 | 1 | 1 | 1 | 1 | 4 | 1 | 3 | 2 | 3 | 1 | 2 | 1 | 1 | 2 | t | 1 | 2 | | | | |
| V ↑ | 1 | v | 1 | 1 | 2 | 4 | 1 | r | b ₁ | b ₂ | b ₃ | t' | 3 | t | t' | 3 | 1 | 1 | 2 | | | | |
| | 1 | v' | 1 | | | 2 | 1 | 1 | 3 | 2 | 3 | 1 | 2 | 1 | 2 | 1 | 3 | 2 | 1 | | | | |
| | 1 | 1 | 1 | | | 1 | 2 | 2 | 1 | | 1 | 1 | 1 | | 1 | 1 | 1 | | | | | | |
| | ⋮ | ⋮ | ⋮ | | | | | | | | | | | | | | | | | | | | |

What about OR?

http://www.claymath.org/Popular_Lectures/Minesweeper/